

Computing Architectures and Algorithms

**Optimizing Gaussian Elimination** 

### **Recall from previous lesson**

- Manufacturers/vendors implement BLAS that are optimized for their machines
- To write optimal code, best to make use of BLAS if possible
  - o BLAS level 3 is the most efficient
  - Better to reorder algorithms to use BLAS 3 versus BLAS level 1 or 2

### **Objective of this lesson**

- Show how an algo can be reordered to make use of BLAS 3
- Reordering uses blocking, so that we operate more on submatrices instead of vectors/scalars

## Reordering Gaussian Elimination to use BLAS3

Primary reference: *Applied Numerical Linear Algebra* by James W. Demmel

 Use Gaussian Elimination which can be defined as:
 "Take each row and subtract multiplies of it from later rows to zero out the entries below the diagonal"

- Use an "in-place" algorithm, where L and U are overwritten on A:
- $\circ\;$  Example, for the given below

$$A = \begin{bmatrix} 2 & 2 & 3 \\ 5 & 9 & 10 \\ 4 & 1 & 2 \end{bmatrix}$$
$$A = LU$$

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 5/2 & 1 & 0 \\ 2 & -3/4 & 1 \end{bmatrix} \quad U = \begin{bmatrix} 2 & 2 & 3 \\ 0 & 4 & 5/2 \\ 0 & 0 & -17/8 \end{bmatrix}$$

- Use an "in-place" algorithm, where L and U are overwritten on A:
- Example, algorithm rewrites A with L and U as shown below:

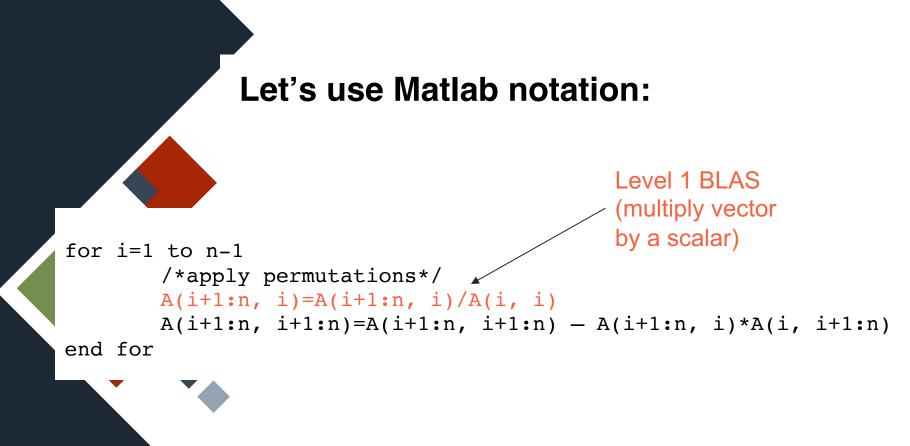
$$A = \begin{bmatrix} 2 & 2 & 3 \\ 5/2 & 4 & 5/2 \\ 2 & -3/4 & -17/8 \end{bmatrix}$$
$$\begin{bmatrix} 1 & 0 & 0 \\ 5/2 & 1 & 0 \\ 2 & -3/4 & 1 \end{bmatrix} \quad U = \begin{bmatrix} 2 & 2 & 3 \\ 0 & 4 & 5/2 \\ 0 & 0 & -17/8 \end{bmatrix}$$

L =

for i=1 to n-1/\* apply permutations so  $a_{ii} \neq 0$  \*/ for j = i + 1 to n $a_{ii} = a_{ii}/a_{ii}$ end for for j = i + 1 to nfor k = i + 1 to n $a_{ik} = a_{ik} - a_{ii} * a_{ik}$ end for end for end for

for i=1 to n-1/\* apply permutations so  $a_{ii} \neq 0$  \*/ for j = i + 1 to n $a_{ii} = a_{ii}/a_{ii}$ end for for j = i + 1 to nfor k = i + 1 to n"Take each row and  $a_{jk} = a_{jk} - a_{ji} * a_{ik}$ subtract multiplies of it from later rows to zero end for out the entries below end for the diagonal" end for

#### Let's use Matlab notation:

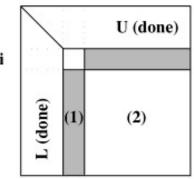


#### Let's use Matlab notation:

```
for i=1 to n-1
        /*apply permutations*/
        A(i+1:n, i) = A(i+1:n, i) / A(i, i)
        A(i+1:n, i+1:n) = A(i+1:n, i+1:n) - A(i+1:n, i) * A(i, i+1:n)
end for
                         Level 2 BLAS
                         (rank-1 update of the submatrix
                         A(i + 1: n, i + 1: n))
```

# We shall reorder the algo to use Level 3 BLAS

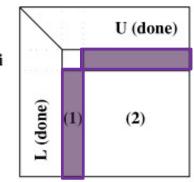
 Modify algorithm slightly to be used within our Level 3 version
 Matrix is now m-by-n



#### Step i of Level 2 BLAS Implementation of LU

# We shall reorder the algo to use Level 3 BLAS

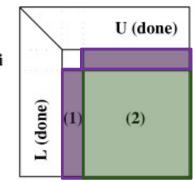
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#### Step i of Level 2 BLAS Implementation of LU

# We shall reorder the algo to use Level 3 BLAS

 Modify algorithm slightly to be used within our Level 3 version
 Matrix is now m-by-n



#### Step i of Level 2 BLAS Implementation of LU

### **Our BLAS 3 version will use blocking**

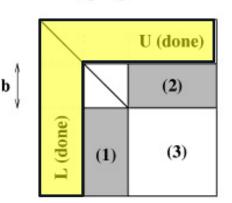
- We will "delay" the update of submatrix 2 by b steps, where b is our block size
- Apply *b* rank-1 updates all at once in one matrix-matrix multiplication

# First, let's see what happens mathematically

 Suppose we are done computing first *i* - 1 columns of *L* and rows of *U*

 $A = \begin{bmatrix} L_{11} & 0 & 0 \\ L_{11} & I & 0 \\ L_{11} & 0 & I \end{bmatrix} \cdot \begin{bmatrix} U_{11} & U_{21} & U_{31} \\ 0 & \tilde{A}_{22} & \tilde{A}_{23} \\ 0 & \tilde{A}_{22} & \tilde{A}_{23} \end{bmatrix}$ 

$$A = \begin{pmatrix} A_{11} & A_{12} & A_{13} \\ A_{21} & A_{22} & A_{23} \\ A_{31} & A_{32} & A_{33} \end{pmatrix}_{n-b-i+1}^{i-1}$$

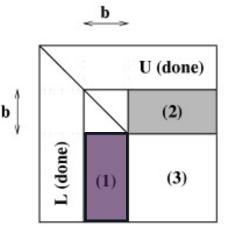


# First, let's see what happens mathematically

• Apply our BLAS2 algorithm to submatrix  $\begin{bmatrix} \tilde{A}_{22} \\ \tilde{A}_{32} \end{bmatrix}$  to get:

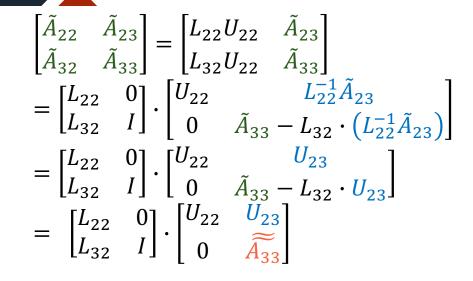
$$\begin{bmatrix} \tilde{A}_{22} \\ \tilde{A}_{32} \end{bmatrix} = \begin{bmatrix} L_{22} \\ L_{32} \end{bmatrix} \cdot U_{22} = \begin{bmatrix} L_{22} U_{22} \\ L_{32} U_{22} \end{bmatrix}$$

We can then write:  $\begin{bmatrix} \tilde{A}_{22} & \tilde{A}_{23} \\ \tilde{A}_{32} & \tilde{A}_{33} \end{bmatrix} = \begin{bmatrix} L_{22}U_{22} & \tilde{A}_{23} \\ L_{32}U_{22} & \tilde{A}_{33} \end{bmatrix}$ 



Step i of Level 3 BLAS Implementation of LU

## First, let's see what happens mathematically

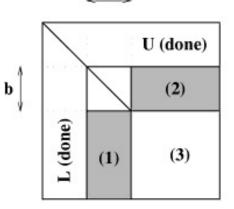


b U (done) (2) (1) (3)

#### Step i of Level 3 BLAS Implementation of LU

## We get updated factorization with b more columns of L and U completed

$$= \begin{bmatrix} L_{22} & 0 \\ L_{32} & I \end{bmatrix} \cdot \begin{bmatrix} U_{22} & U_{23} \\ 0 & \widetilde{A_{33}} \end{bmatrix}$$
$$\begin{bmatrix} A_{11} & A_{12} & A_{13} \\ A_{21} & A_{22} & A_{23} \\ A_{31} & A_{32} & A_{33} \end{bmatrix} = \begin{bmatrix} L_{11} & 0 & 0 \\ L_{21} & L_{22} & 0 \\ L_{31} & L_{23} & I \end{bmatrix} \cdot \begin{bmatrix} U_{11} & U_{21} & U_{31} \\ 0 & U_{22} & U_{23} \\ 0 & 0 & \widetilde{A_{33}} \end{bmatrix}$$



b

#### Step i of Level 3 BLAS Implementation of LU

### **BLAS L3 Algorithm for LU Factorization**

1) Use BLAS L2 Algorithm to factorize  $\begin{bmatrix} \tilde{A}_{22} \\ \tilde{A}_{32} \end{bmatrix} = \begin{bmatrix} L_{22} \\ L_{32} \end{bmatrix} \cdot U_{22}$ 2) Form  $U_{23} = L_{22}^{-1}\tilde{A}_{23}$  (this is a BLAS L3 operation) 3) Form  $\widetilde{\tilde{A}_{33}} = \tilde{A}_{33} - L_{32} \cdot (L_{22}^{-1}\tilde{A}_{23})$  (MMM operation, BLAS L3)

### **BLAS L3 Algorithm for LU Factorization**

for i = 1 to n - 1 step bFactorize  $A(i:n, i:i + b) = \begin{bmatrix} L_{22} \\ L_{32} \end{bmatrix} U_{22}$ /\* use BLAS L2 algo\*/  $A(i:i + b - 1, i + b:n) = L_{22}^{-1} \cdot A(i:i + b - 1, i + b:n)$ /\* Form  $U_{23}$  \*/  $A(i + b:n, i + b:n) = A(i + b:n, i + b:n) - A(i + b:n, i:i + b - 1) \cdot A(i:i + b - 1, i + b:n)$ /\* Form  $\widetilde{A_{33}}$  \*/

end for

### **Additional remarks**

- Need to choose block size to maximize speed Ο • Large blocks to multiply larger matrices Number of floating point operations by Level 2 and Level 1 BLAS in step 1 is about  $\frac{n^2b}{2}$  for small b Grows as b grows, we don't want to pick b too large Commonly used values are b = 32 or b = 64Detailed implementations: BLAS 2 Algo: sgetf2 on LAPACK BLAS 3 Algo: sgetrf on LAPACK Ο
  - Search here: http://www.netlib.org/lapack/explorehtml/modules.html